

NonLocal Boxes
& *Communication Complexity*

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(Tuesday 13th of September, 2022.)

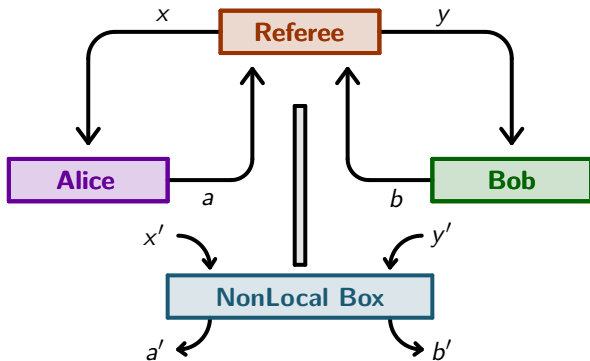
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— *Part 1* —

Definitions & Notations

CHSH Game

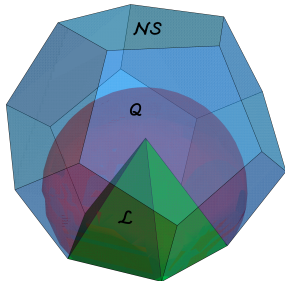


Win at CHSH. $a \oplus b = x y$.

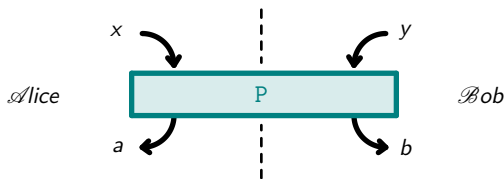
Win at CHSH'. $a \oplus b = (x \oplus 1)(y \oplus 1)$.

Strategies

- **Deterministic strategies.** $\rightsquigarrow \max \mathbb{P}(\text{win}) = 75\%$.
- **Classical strategies \mathcal{L} .** $\rightsquigarrow \max \mathbb{P}(\text{win}) = 75\%$.
- **Quantum strategies \mathcal{Q} .** $\rightsquigarrow \max \mathbb{P}(\text{win}) = \cos^2\left(\frac{\pi}{8}\right) \approx 85\%$.
- **Non-signalling strategies \mathcal{NS} .** $\rightsquigarrow \max \mathbb{P}(\text{win}) = 100\%$.



NonLocal Box

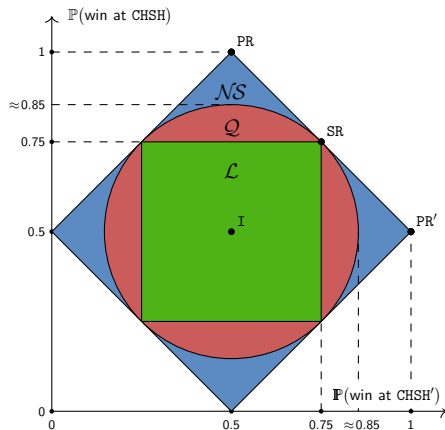


Definition. A **nonlocal box** is a function:

$$P : \begin{cases} \{0, 1\}^4 & \longrightarrow [0, 1] \\ (a, b, x, y) & \longmapsto P(a, b \mid x, y). \end{cases}$$

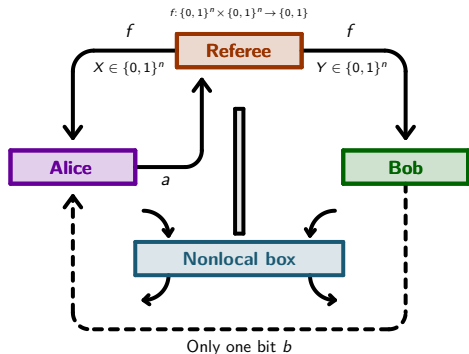
such that (i) P is a conditional probability distribution and (ii) $P \in \mathcal{NS} \setminus \mathcal{L}$.

Examples



- $\text{PR}(a, b | x, y) := \begin{cases} \frac{1}{2} & \text{si } a \oplus b = xy, \\ 0 & \text{otherwise.} \end{cases}$
- $\text{PR}'(a, b | x, y) := \begin{cases} \frac{1}{2} & \text{si } a \oplus b = (x \oplus 1)(y \oplus 1), \\ 0 & \text{otherwise.} \end{cases}$
- $\text{SR}(a, b | x, y) := \begin{cases} \frac{1}{2} & \text{si } a = b, \\ 0 & \text{otherwise.} \end{cases}$
- $\text{I}(a, b | x, y) := \frac{1}{4}$

Communication Complexity



$$\text{Win} \iff a = f(X, Y).$$

Def. A function f is said to be **trivial** (in the sense of communication complexity) if Alice knows any value $f(X, Y)$ with only one bit transmitted between Alice and Bob.

Ex. For $n = 2$, $X = (x_1, x_2)$, $Y = (y_1, y_2)$:

- $f := x_1 \oplus y_1 \oplus x_2 \oplus y_2 \oplus 1$ is trivial.
- $g := (x_1 x_2) \oplus (y_1 y_2)$ is trivial.
- $h := (x_1 y_1) \oplus (x_2 y_2)$ is NOT trivial.

Def. A box P is said to be **collapsing** (or trivial) if using copies of this box P any Boolean function f is trivial, with probability $\geq q > \frac{1}{2}$.

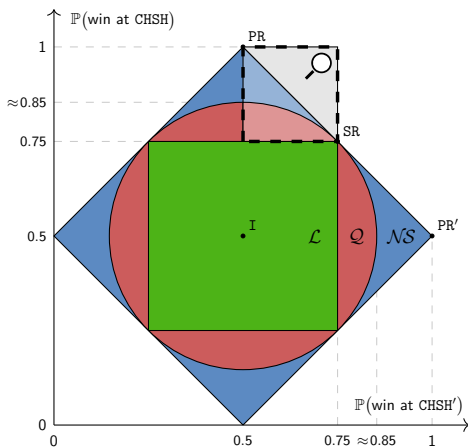
Ex. Link with previous boxes:

- The boxes PR and PR' are collapsing.
- The boxes SR and I are NOT collapsing.

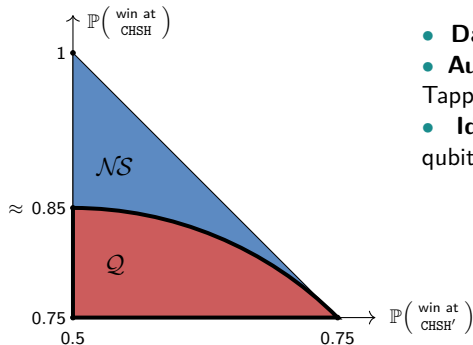
— *Part 2* —

Historical Overview

Goal. Show that quantum boxes are **non-collapsing** but that post-quantum boxes are **collapsing**.

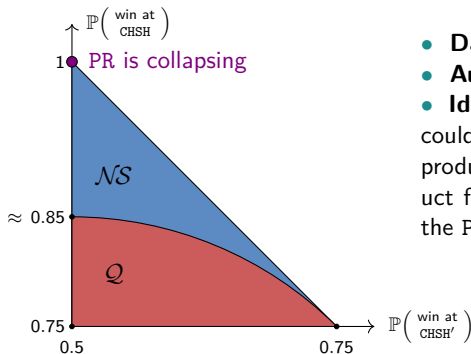


1999: Quantum boxes are non-collapsing



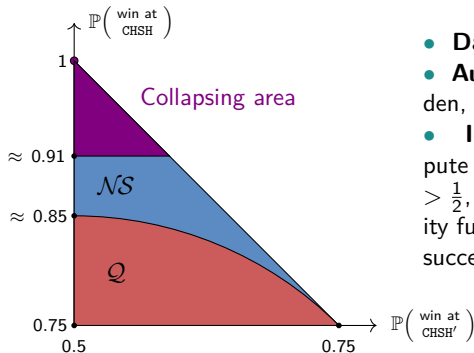
- **Date.** 1999 [1].
- **Authors.** Cleve, van Dam, Nielsen, Tapp.
- **Ideas.** (1) Prove the result with qubits, (2) Go back to bits.

1999: The PR box is collapsing



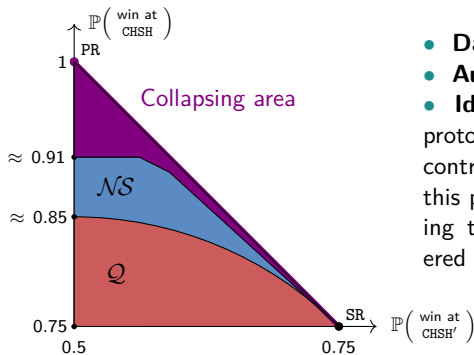
- **Date.** 1999 [2].
- **Author.** van Dam.
- **Ideas.** (1) Any Boolean function f could be written in terms of an inner product function, (2) Any inner product function is trivial using copies of the PR box.

2006: Boxes above $\approx 91\%$ are collapsing



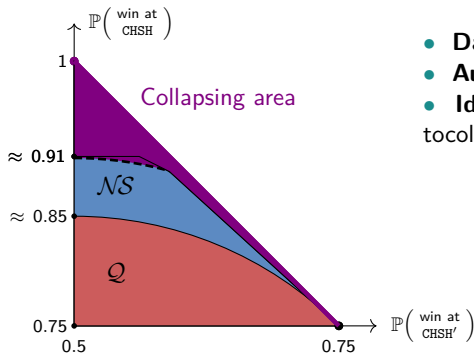
- **Date.** 2006 [3].
- **Authors.** Brassard, Buhrman, Linden, Méthot, Tapp, Unger.
- **Ideas.** (1) Distributively compute the given function f with proba $> \frac{1}{2}$, (2) Inductively apply the majority function Maj in order to boost the success probability.

2009: Correlated boxes are collapsing



- **Date.** 2009 [4].
- **Authors.** Brunner, Skrzypczyk.
- **Ideas.** (1) Introduce a distillation protocol, cf. generalization in "Our contribution", (2) Inductively apply this protocol many times until reaching the "collapsing triangle" discovered in 2006.

2018: Boxes above an ellipse are collapsing

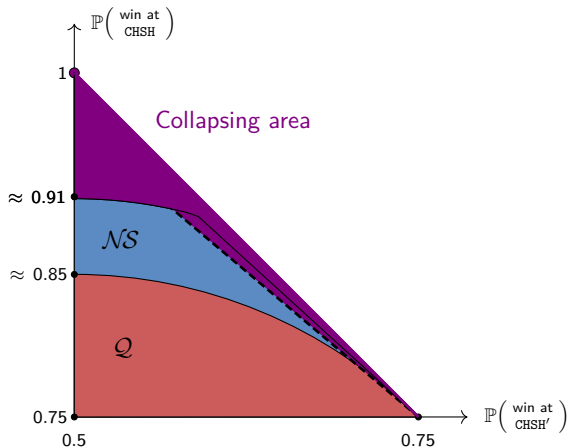


- **Date.** 2018 [5].
- **Author.** Broadbent, Proulx.
- **Idea.** Generalize BBLMTU's protocol (cf. 2006).

— *Part 3* —

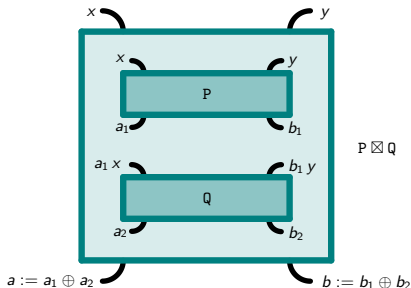
**Our Contribution:
Algebra of Boxes**

Our contribution



Algebra of boxes

Recall. A nonlocal box P is a conditional probability distribution $(a, b, x, y) \in \{0, 1\}^4 \mapsto P(a, b | x, y) \in [0, 1]$ such that $P \in \mathcal{NS} \setminus \mathcal{L}$.



$$P \otimes Q(a, b | x, y) := \sum_{a_1, b_1 \in \{0, 1\}} P(a_1, b_1 | x, y) \times Q(a \oplus a_1, b \oplus b_1 | a_1 x, b_1 y)$$

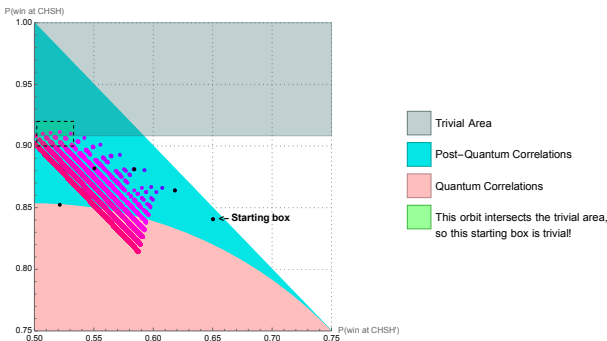
Algebra of boxes. The vector space $\mathcal{B} := \mathcal{F}(\{0, 1\}^4, \mathbb{R})$ endowed with the operations $\{+, \cdot, \otimes\}$ defines a non-commutative and non-associative algebra.

Orbit of a box

Orbit of order k . $\text{Orbit}_k(P) := \{ \text{products of exactly } k \text{ times the term } P \}$.

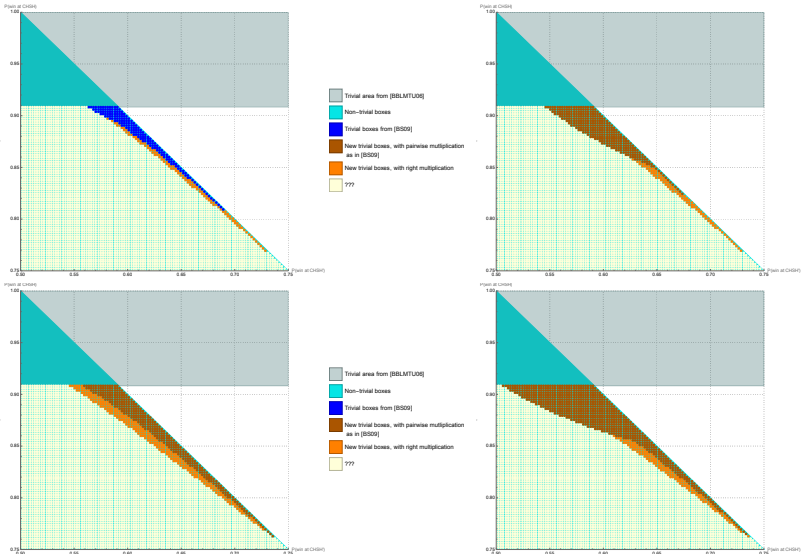
Examples. • $\text{Orbit}_3(P) = \{ P \boxtimes (P \boxtimes P), (P \boxtimes P) \boxtimes P \},$

• $\text{Orbit}_4(P) = \{ P \boxtimes (P \boxtimes (P \boxtimes P)), P \boxtimes ((P \boxtimes P) \boxtimes P), (P \boxtimes (P \boxtimes P)) \boxtimes P, ((P \boxtimes P) \boxtimes P) \boxtimes P, (P \boxtimes P) \boxtimes (P \boxtimes P) \}.$



The "highest" box in each orbit. $P_{\max, k} = (((P \boxtimes P) \boxtimes P) \dots) \boxtimes P =: P^{\boxtimes k}.$

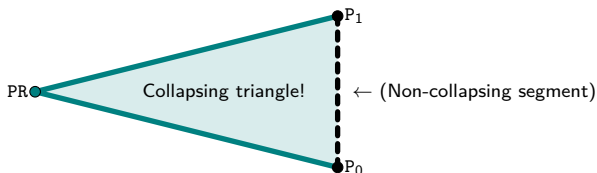
New collapsing boxes: numerically



New collapsing boxes: analytically

Theorem 1 (New collapsing boxes)

In the triangle whose vertices are $\{PR, P_0 := \mathbf{1}_{a=b=0}, P_1 := \mathbf{1}_{a=b=1}\}$, all the points are collapsing boxes, except points in the segment P_0 - P_1 .



Proof. (1) The triangle is stable under \boxtimes . (2) Define a sequence: initialize at an arbitrary point of the triangle (except in the vertical segment), and inductively apply the multiplication \boxtimes . (3) This sequence converges to PR. (4) But, near PR, all boxes are collapsing (cf. 2006). (5) Hence, the orbit intersects the collapsing area and the starting box must be collapsing as well. \square

Bibliography

- [1] R. Cleve, W. van Dam, M. Nielsen, and A. Tapp, *Quantum Entanglement and the Communication Complexity of the Inner Product Function*. Berlin, Heidelberg: Springer Berlin Heidelberg, 1999.
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- [5] M.-O. Proulx, "A limit on quantum nonlocality from an information processing principle," Master's thesis, Department of Physics, University of Ottawa, Canada, 2018.
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Under the supervision of Anne Broadbent, Ion Nechita and Clément Pellegrini.